

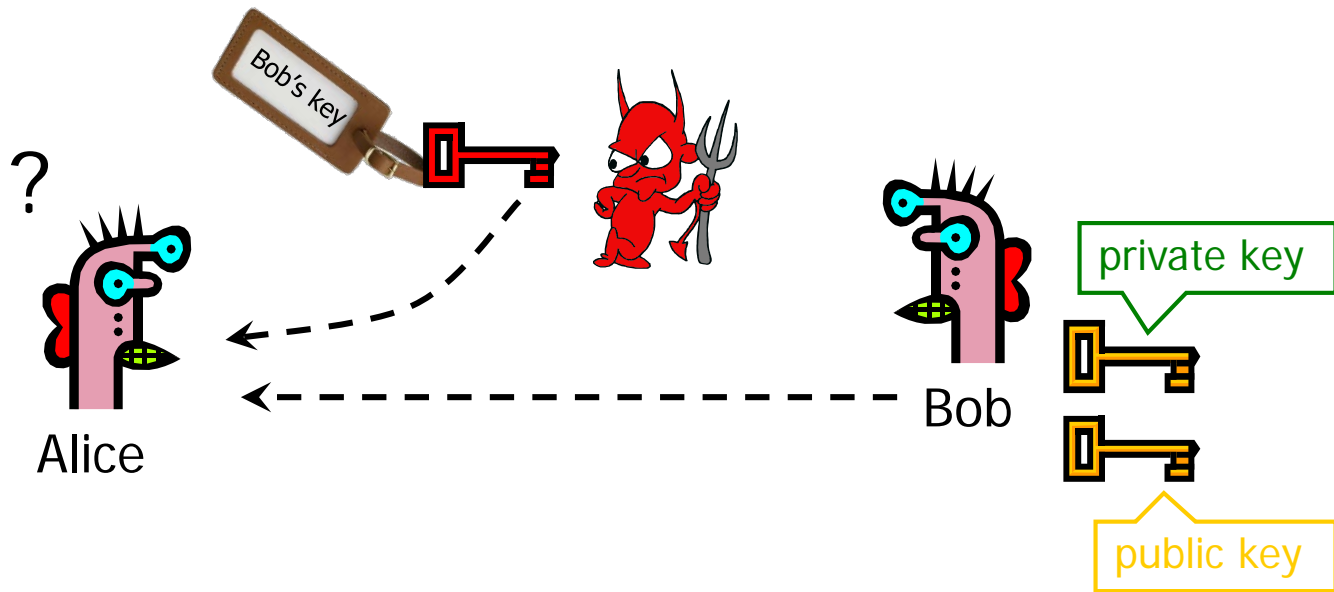
# Public-Key Infrastructure

## NETS E2008

---

Many slides from  
Vitaly Shmatikov, UT Austin

# Authenticity of Public Keys



Problem: How does Alice know that the public key she received is really Bob's public key?

# Distribution of Public Keys

---

- ◆ Public announcement or public directory
  - Risks: forgery and tampering
- ◆ Public-key certificate
  - Signed statement specifying the key and identity
    - $\text{sig}_{\text{Alice}}(\text{"Bob"}, \text{PK}_B)$
- ◆ Common approach: certificate authority (CA)
  - Single agency responsible for certifying public keys
  - After generating a private/public key pair, user proves his identity and knowledge of the private key to obtain CA's certificate for the public key (offline)
  - Every computer is pre-configured with CA's public key

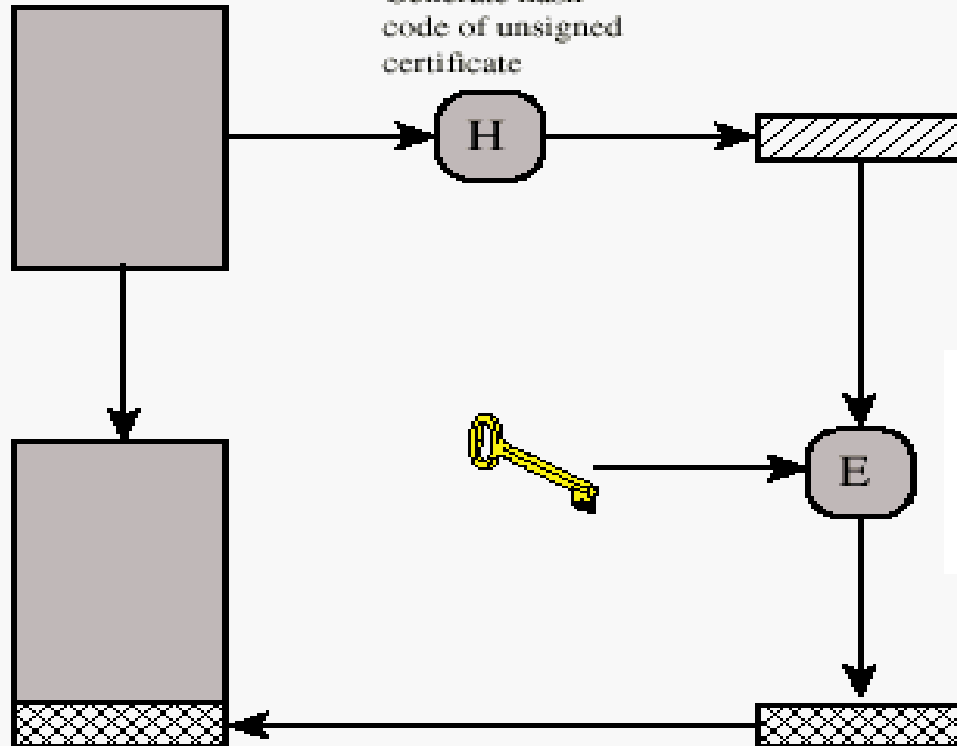
# Obtaining a User's Certificate

---

- ◆ Characteristics of certificates generated by CA:
  - Any user with access to the public key of the CA can verify the user public key that was certified.
  - No part other than the CA can modify the certificate without this being detected.

# Using Public-Key Certificates

Unsigned certificate:  
contains user ID,  
user's public key



Authenticity of public keys is reduced to  
authenticity of one key (CA's public key)

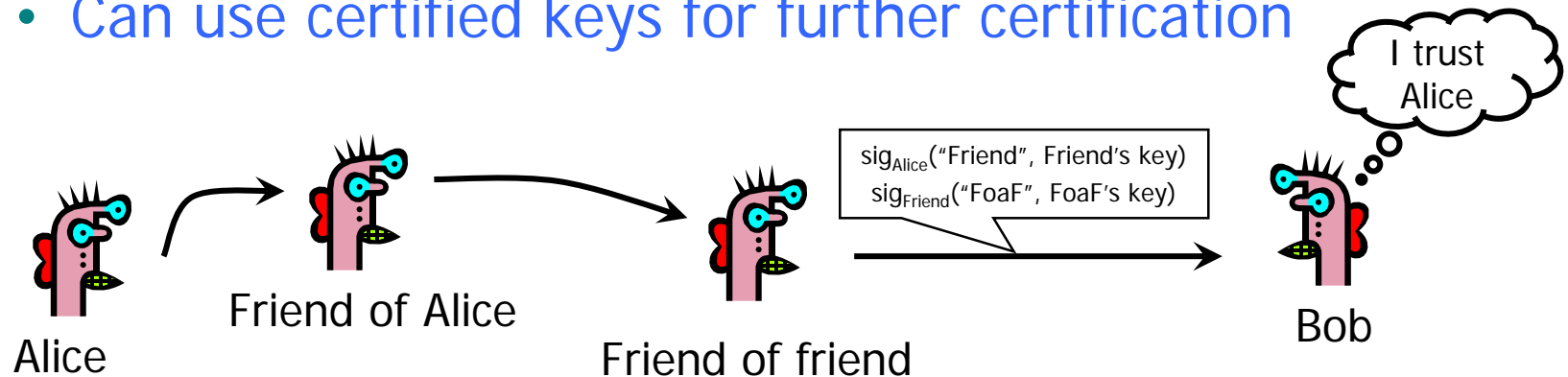
# Hierarchical Approach

---

- ◆ Single CA certifying every public key is impractical
- ◆ Instead, use a trusted **root authority**
  - For example, Verisign
  - Everybody must know the public key for verifying root authority's signatures
- ◆ Root authority signs certificates for lower-level authorities, lower-level authorities sign certificates for individual networks, and so on
  - Instead of a single certificate, use a **certificate chain**
    - $\text{sig}_{\text{Verisign}}(\text{"UT Austin"}, \text{PK}_{\text{UT}})$ ,  $\text{sig}_{\text{UT}}(\text{"Vitaly S."}, \text{PK}_V)$
  - What happens if root authority is ever compromised?

# Alternative: "Web of Trust"

- ◆ Used in PGP (Pretty Good Privacy)
- ◆ Instead of a single root certificate authority, each person has a set of keys they "trust"
  - If public-key certificate is signed by one of the "trusted" keys, the public key contained in it will be deemed valid
- ◆ Trust can be transitive
  - Can use certified keys for further certification



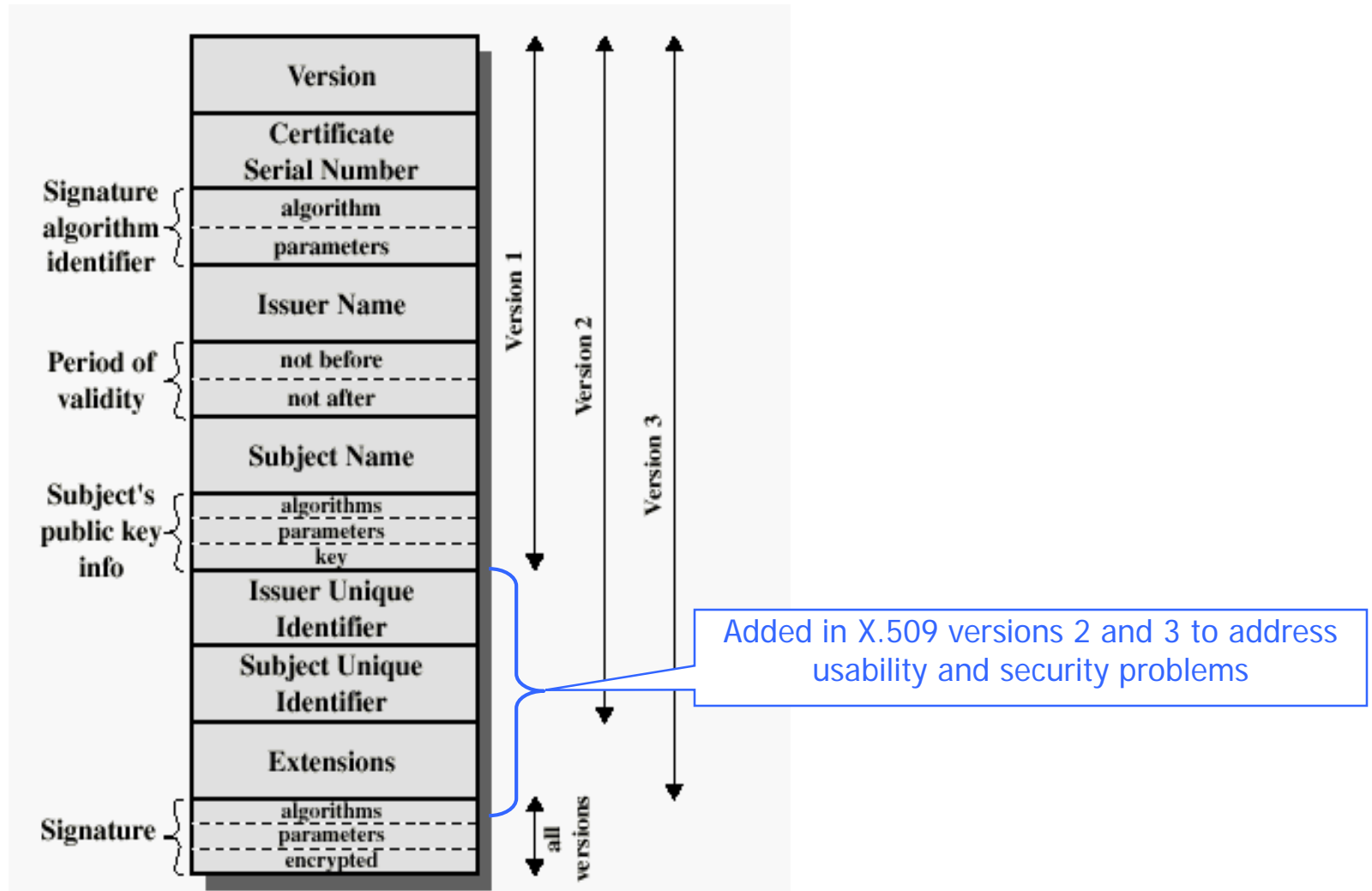
# X.509 Authentication Service

---

- ◆ Internet standard (1988-2000)
- ◆ Specifies certificate format
  - X.509 certificates are used in IPsec and SSL/TLS
- ◆ Specifies certificate directory service
  - For retrieving other users' CA-certified public keys
- ◆ Specifies a set of authentication protocols
  - For proving identity using public-key signatures
- ◆ Does not specify crypto algorithms
  - Can use it with any digital signature scheme and hash function, but hashing is required before signing



# X.509 Certificate



# Certificate Revocation

---

- ◆ Revocation is very important
- ◆ Many valid reasons to revoke a certificate
  - Private key corresponding to the certified public key has been compromised
  - User stopped paying his certification fee to this CA and CA no longer wishes to certify him
  - CA's certificate has been compromised!
- ◆ Expiration is a form of revocation, too
  - Many deployed systems don't bother with revocation
  - Re-issuance of certificates is a big revenue source for certificate authorities

# Certificate Revocation Mechanisms

---

## ◆ Online revocation service

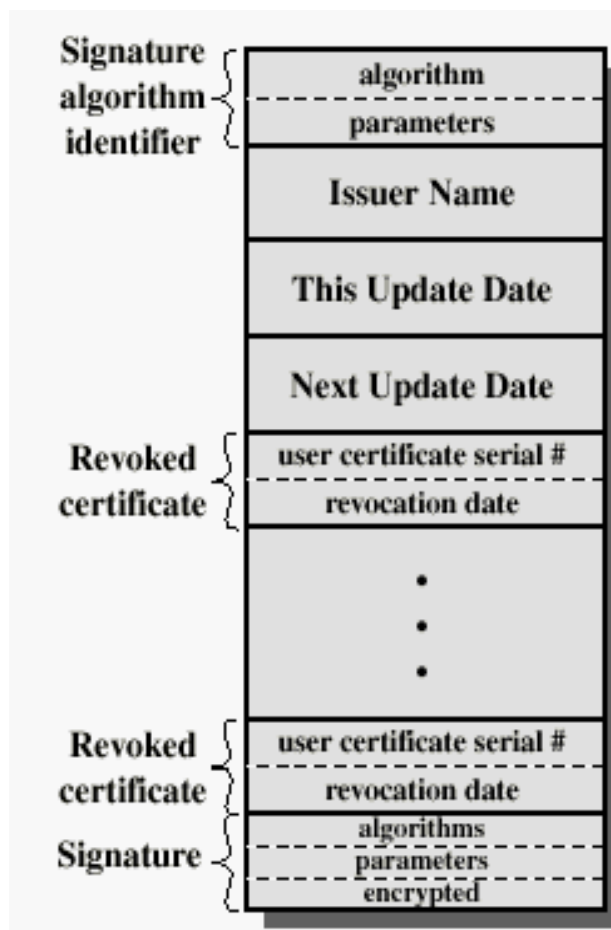
- When a certificate is presented, recipient goes to a special online service to verify whether it is still valid
  - Like a merchant dialing up the credit card processor

## ◆ Certificate revocation list (CRL)

- CA periodically issues a signed list of revoked certificates
  - Credit card companies used to issue thick books of canceled credit card numbers
- Can issue a “delta CRL” containing only updates

## ◆ Question: does revocation protect against forged certificates?

# X.509 Certificate Revocation List



Because certificate serial numbers must be unique within each CA, this is enough to identify the certificate

# Online Certificate Status Protocol

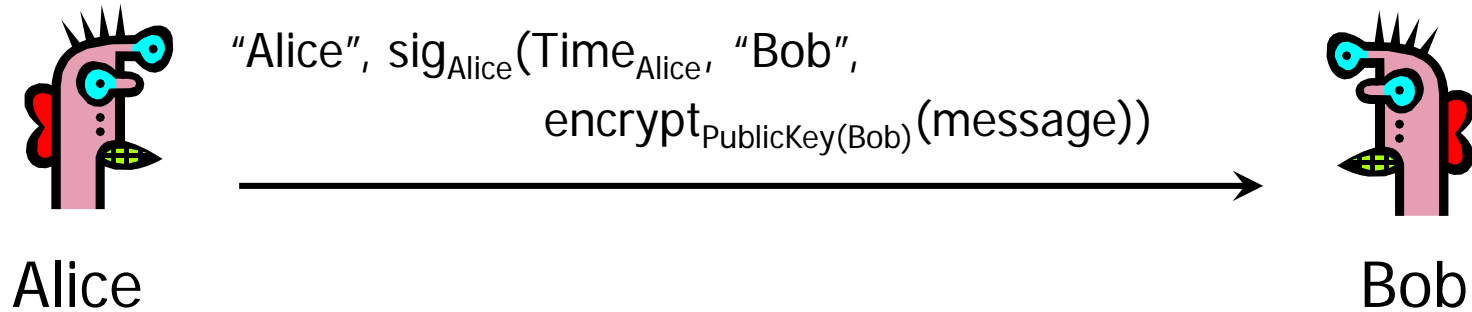
---

## ◆ RFC 2560

- Saves retrieving the complete CRL
- OCSP responders could be chained to some degree
  - eg. trusted responder could query other CA's OCSP

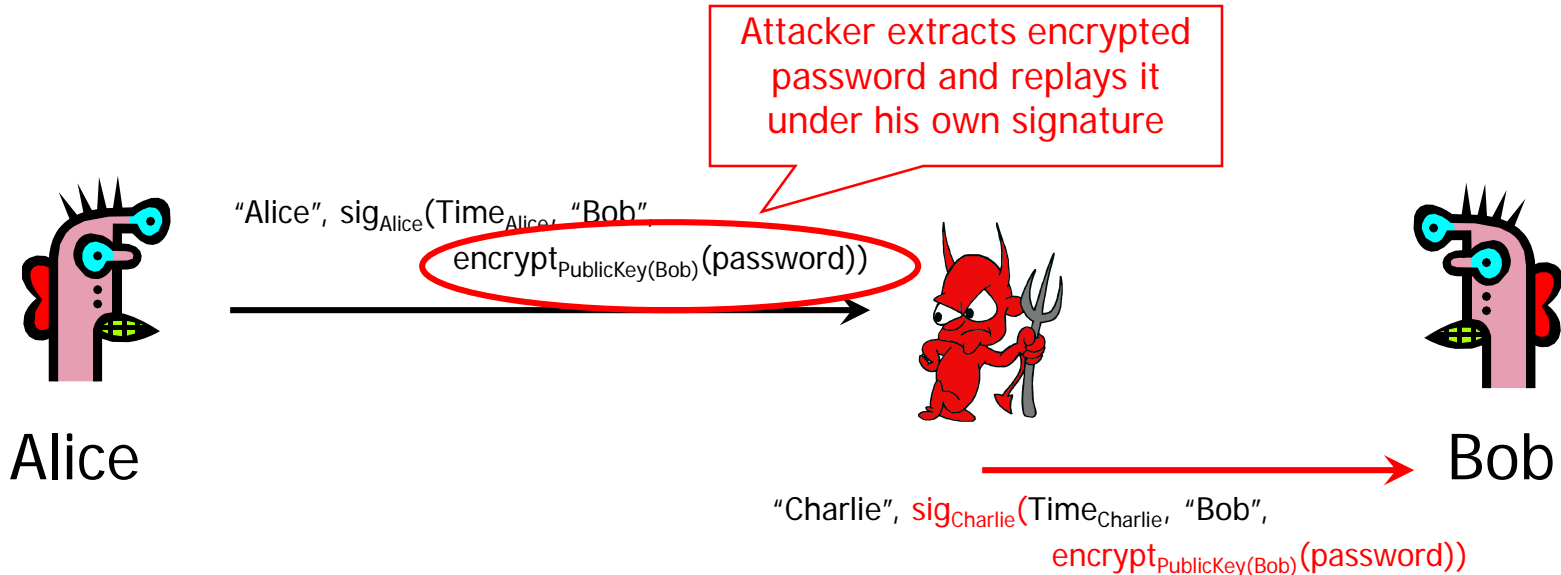
# X.509 Version 1

---



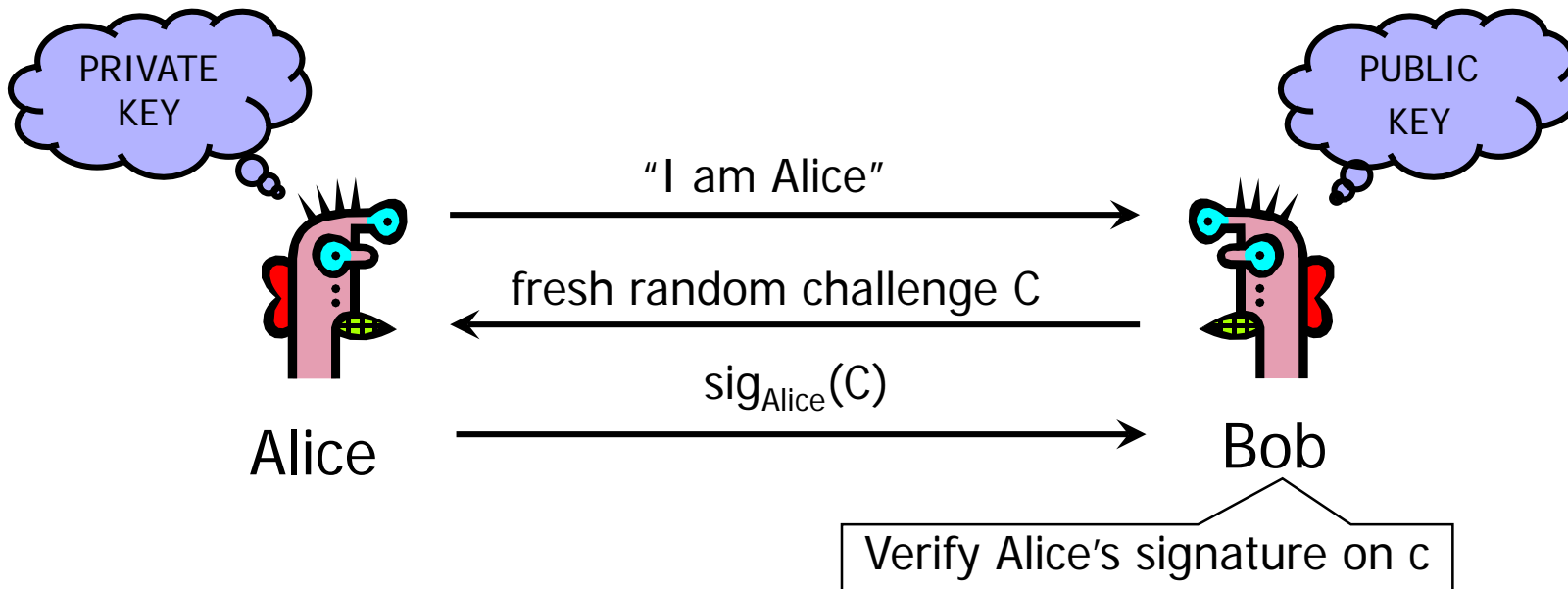
- ◆ Encrypt, then sign for **authenticated encryption**
  - Goal: achieve both confidentiality and authentication
  - E.g., encrypted, signed password for access control
- ◆ Does this work?

# Attack on X.509 Version 1



- ◆ Receiving encrypted password under signature does not mean that the sender actually knows the password!
- ◆ Proper usage: **sign, then encrypt**

# Authentication with Public Keys



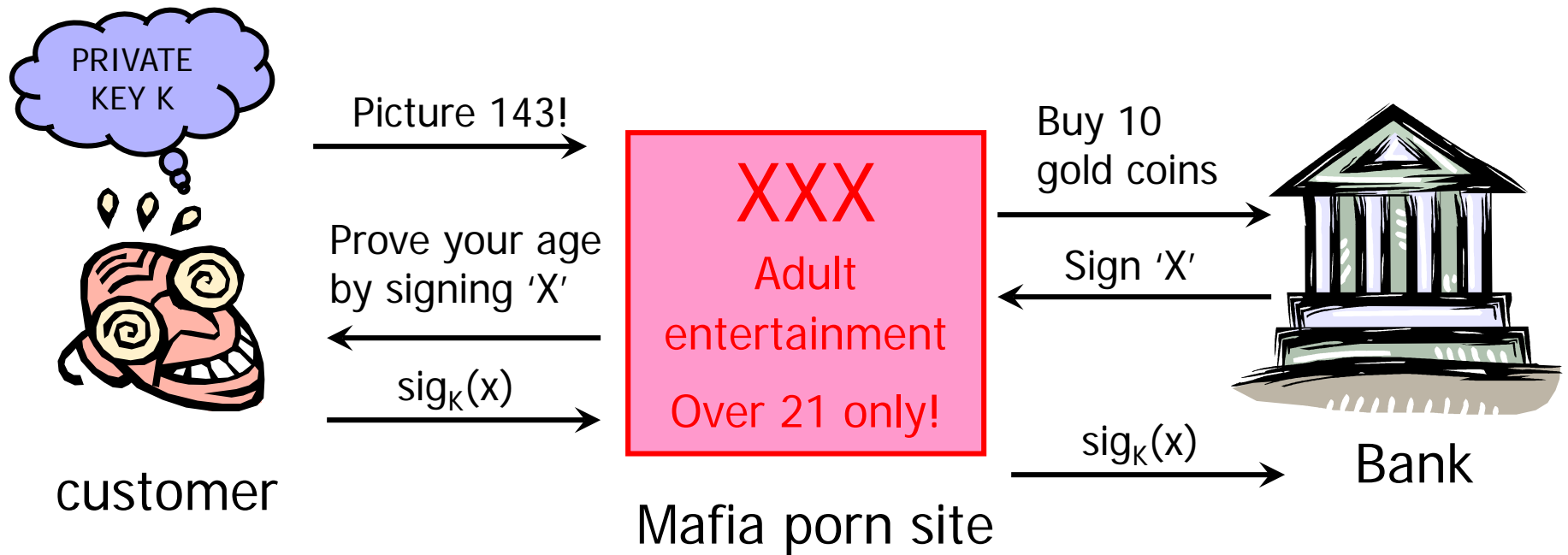
1. Only Alice can create a valid signature
2. Signature is on a fresh, unpredictable challenge

Potential problem: Alice will sign anything

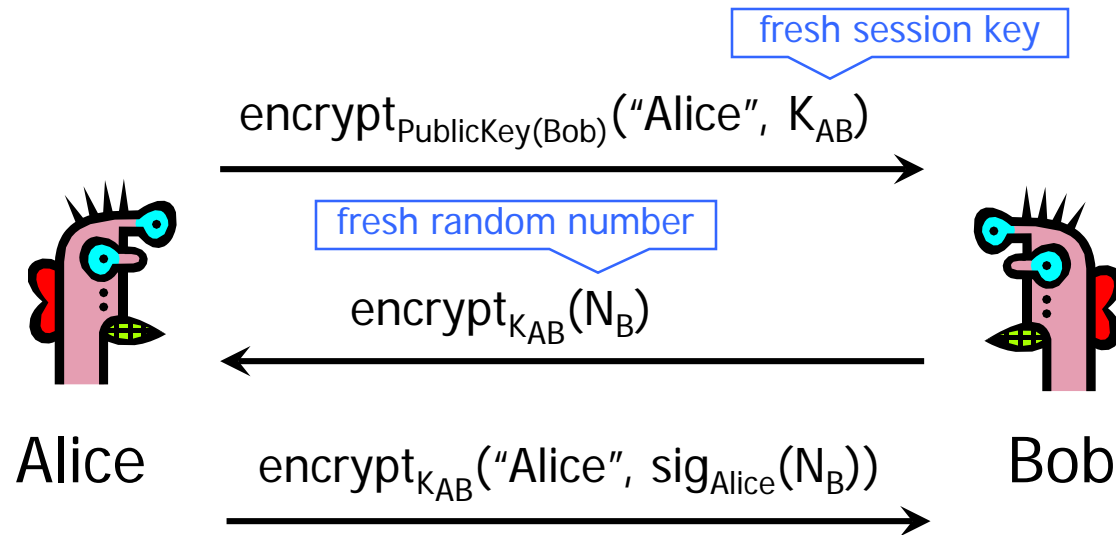


# Mafia-in-the-Middle Attack

[from Anderson's book]



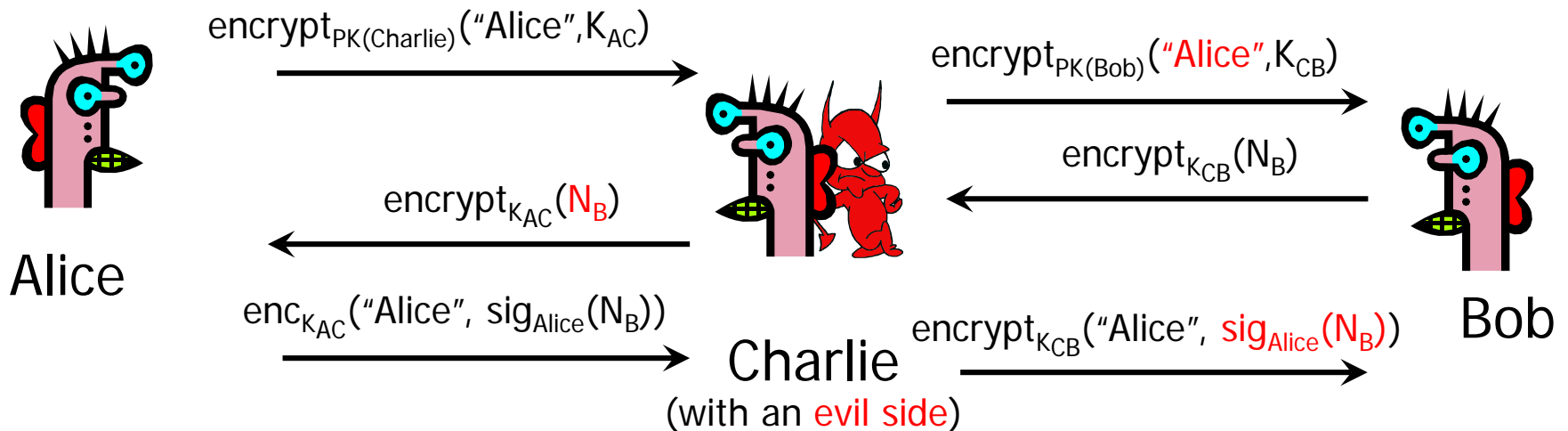
# Early Version of SSL (Simplified)



## ◆ Bob's reasoning: I must be talking to Alice because...

- Whoever signed  $N_B$  knows Alice's private key... Only Alice knows her private key... Alice must have signed  $N_B$ ...  $N_B$  is fresh and random and I sent it encrypted under  $K_{AB}$ ... Alice could have learned  $N_B$  only if she knows  $K_{AB}$ ... She must be the person who sent me  $K_{AB}$  in the first message...

# Breaking Early SSL



- ◆ Charlie uses his legitimate conversation with Alice to impersonate Alice to Bob
  - Information signed by Alice is not sufficiently explicit

# More Litterature

---

- ◆ Wikipedia entry on X.509
  - [Contains list of different file formats](#)
- ◆ RFC 3280 "Internet X.509 Public Key Infrastructure Certificate and Certificate Revocation List (CRL) Profile"
- ◆ IETF PKIX charter
  - <http://www.ietf.org/html.charters/pkix-charter.html>
- ◆ [www.openvalidation.org](http://www.openvalidation.org)
  - OCSP validation resources
- ◆ [www.openca.org](http://www.openca.org)
  - Open Source CA and OCSP software